

## SOME PROPERTIES ON DERIVATIONS OF LATTICES

M. F. KAWAGUCHI AND M. KONDO\*

ABSTRACT. In this paper we consider some properties of derivations of lattices and show that (i) for a derivation  $d$  of a lattice  $L$  with the maximum element  $1$ , it is monotone if and only if  $d(x) \leq d(1)$  for all  $x \in L$  (ii) a monotone derivation  $d$  is characterized by  $d(x) = x \wedge d(1)$  and (iii) simple characterization theorems of modular lattices and of distributive lattices are given by derivations. We also show that, for a distributive lattice  $L$  and a monotone derivation  $d$  of it, the set  $\text{Fix}_d(L)$  of all fixed points of  $d$  is isomorphic to the lattice  $L/\ker(d)$ . We provide a counter example to the result (Theorem 4) proved in [3].

### 1. INTRODUCTION

A notion of derivations of algebras with two operations  $+$  and  $\cdot$  has introduced as an analogy of derivations of analysis and then some properties of derivations are considered. For an algebra  $A = (A, +, \cdot)$ , a map  $f : A \rightarrow A$  is called a derivation if it satisfies the conditions, for all  $x, y \in A$ ,

$$\begin{aligned}f(x + y) &= f(x) + f(y) \\f(x \cdot y) &= f(x) \cdot y + x \cdot f(y).\end{aligned}$$

The notion of derivation is important in the theory of rings ([5]). After that, it is applied to lattices ([4]), where operation  $+$  and  $\cdot$  are interpreted as lattice operations  $\vee$  and  $\wedge$ , respectively. Following the naive

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\*Corresponding author.

interpretation, the derivation  $d$  of a lattice  $L$  may be defined by

$$(a) \quad d(x \vee y) = d(x) \vee d(y)$$

$$(b) \quad d(x \wedge y) = (d(x) \wedge y) \vee (x \wedge d(y)).$$

As proved in [4, 6], the condition (a) represents monotonicity of  $d$  and the condition (b) is equivalent to the condition  $d(x \wedge y) = d(x) \wedge y$ . Hence, as proved later, a monotone derivation  $f : L \rightarrow L$  is characterized by  $f(x \wedge y) = f(x) \wedge y$  for all  $x, y \in L$ . It follows from the result that a monotone derivation  $d$  has the form of  $d(x) = x \wedge d(1)$  if  $L$  has the maximum element 1 and thus every monotone derivation is determined completely by the value  $d(1)$ .

In order to obtain more interesting properties of derivations of lattices, we adopt another definition of derivations according to [1, 2, 3, 7] and prove some fundamental properties of them, from which we get new results about derivations of lattices and provide accurate statements described in [1, 2, 3, 6, 7]. Moreover, we consider properties of generalized derivation ([1, 2]) and of  $f$ -derivation ([3]) from our view point and give some results which have simpler proofs than those of [3].

Concretely, we prove that

- (i). For a derivation  $d$  of a lattice  $L$  with the maximum element 1, it is monotone if and only if  $d(x) \leq d(1)$  for all  $x \in L$ .
- (ii). A monotone derivation  $d$  is just the form of  $d(x) = x \wedge d(1)$ .
- (iii). For any lattice  $L$  and a derivation  $d$ , the condition

$$d \text{ is monotone} \Leftrightarrow d(d(x) \vee y) = d(x) \vee d(y) \quad (\forall x, y \in L)$$

is equivalent to that  $L$  is a modular lattice.

- (iv). For any lattice  $L$  and a derivation  $d$ , the condition

$$d \text{ is monotone} \Leftrightarrow d(x \vee y) = d(x) \vee d(y) \quad (\forall x, y \in L),$$

is equivalent to that  $L$  is a distributive lattice.

We also show that, for a distributive lattice  $L$  and a monotone derivation  $d$  of it, the set  $\text{Fix}_d(L) = \{x \in L \mid d(x) = x\}$  of all fixed points of  $d$  is isomorphic to the lattice  $L/\ker(d)$ .

Lastly, we provide a counter example to the result (Theorem 4) proved in [3].

## 2. DERIVATION

According to [6, 7], we give a definition of derivation of a lattice. Let  $L = (L, \vee, \wedge)$  be a lattice. A map  $d : L \rightarrow L$  is called a *derivation* of  $L$  if it satisfies the condition

$$d(x \wedge y) = (d(x) \wedge y) \vee (x \wedge d(y)) \quad (\forall x, y \in L)$$

Moreover, a derivation  $d$  is called *monotone* if

$$x \leq y \Rightarrow d(x) \leq d(y) \quad (\forall x, y \in L).$$

We note that the notion of monotone is called isotone in [1, 2, 3, 7])

**Example 1.** Let  $L$  be a lattice and  $a \in L$ . If we define a map  $d_a : L \rightarrow L$  by  $d_a(x) = x \wedge a$ , then  $d_a$  is a monotone derivation. Indeed, for all  $x, y \in L$ , we have  $d_a(x \wedge y) = (x \wedge y) \wedge a = ((x \wedge a) \wedge y) \vee (x \wedge (y \wedge a)) = (d_a(x) \wedge y) \vee (x \wedge d_a(y))$ .

**Example 2.** ([3]) Let  $L = \{0, a, b, 1\}$ , ( $0 < a < b < c < 1$ ). We define  $d : L \rightarrow L$  by

$$d(x) = \begin{cases} 0 & (x = 0) \\ a & (x = a, b) \\ c & (x = c, 1) \end{cases}$$

It is clear that  $d : L \rightarrow L$  is the derivation of  $L$ .

We have basic results about derivations of lattices.

**Proposition 2.1.** *Let  $L$  be a lattice and  $d$  be a derivation of  $L$ . For all  $x, y \in L$ ,*

- (1)  $d(x) \leq x$
- (2)  $d(d(x)) = d(x)$
- (3) *If  $1 \in L$ , then  $d(x) = d(x) \vee (x \wedge d(1))$*
- (4) *If  $1 \in L$ , then  $d(1) = 1 \Leftrightarrow d = id_L$*
- (5)  $d(x) \wedge d(y) \leq d(x \wedge y) \leq d(x) \vee d(y)$
- (6)  $d(d(x) \wedge d(y)) = d(x) \wedge d(y)$
- (7) *If  $d$  is monotone, then  $d(d(x) \vee d(y)) = d(x) \vee d(y)$*
- (8) *If  $d(d(x) \vee y) = d(x) \vee d(y)$ , then  $d$  is monotone.*

*Proof.* We only prove (7) and (8).

(7) If  $d$  is monotone, since  $d(x), d(y) \leq d(x) \vee d(y)$ , then we get  $d(d(x)), d(d(y)) \leq d(d(x) \vee d(y))$ . By  $d(d(x)) = d(x)$  and  $d(d(y)) = d(y)$ , we have  $d(x), d(y) \leq d(d(x) \vee d(y))$  and  $d(x) \vee d(y) \leq d(d(x) \vee d(y))$ . It is clear from (1) that  $d(d(x) \vee d(y)) \leq d(x) \vee d(y)$ . Hence,  $d(d(x) \vee d(y)) = d(x) \vee d(y)$ .

(8) Suppose that  $x \leq y$ . Since  $d(x) \leq x \leq y$ , we have  $d(y) = d(d(x) \vee y) = d(x) \vee d(y)$  and thus  $d(x) \leq d(y)$ .  $\square$

We note that the derivation  $d_a(x) = x \wedge a$  in Example 1 is monotone. Moreover, any monotone derivation  $d$  has just the form of  $d(x) = x \wedge a$  for some  $a \in L$ . In order to prove this fact, we deeply think about properties of monotone derivations.

**Theorem 2.2.** *For any derivation  $d$ , the following conditions are equivalent to each other.*

- (1)  $d$  is monotone ;
- (2)  $d(x \wedge y) = d(x) \wedge d(y) \quad (\forall x, y \in L)$ ;
- (3)  $d(x) \vee d(y) \leq d(x \vee y) \quad (\forall x, y \in L)$ .

*Proof.* We only show the case (1)  $\Rightarrow$  (2). The other cases can be proved easily.

Since  $x \wedge y \leq x, y$ , we have  $d(x \wedge y) \leq d(x), d(y)$ . On the other hand, since  $d(x \wedge y) \leq d(x) \wedge d(y) \leq x \wedge y$ , we get  $d(x \wedge y) = d(d(x \wedge y)) \leq d(d(x) \wedge d(y)) \leq d(x \wedge y)$ . Thus  $d(x \wedge y) = d(d(x) \wedge d(y))$ . It follows that

$$\begin{aligned}
 d(x \wedge y) &= d(d(x) \wedge d(y)) \\
 &= \{d(d(x)) \wedge d(y)\} \vee \{d(x) \wedge d(d(y))\} \\
 &= (d(x) \wedge d(y)) \vee (d(x) \wedge d(y)) \\
 &= d(x) \wedge d(y).
 \end{aligned}$$

□

From the result above, a monotone derivation can be characterized as follows.

**Theorem 2.3.** *Let  $L$  be a lattice and  $f : L \rightarrow L$  be a map. Then the following conditions are equivalent.*

- (1)  $f$  is a monotone derivation;
- (2)  $f(x \wedge y) = f(x) \wedge y \quad (\forall x, y \in L)$ ;
- (3)  $f(x) = x \wedge f(1) \quad (\forall x \in L)$ .

*Proof.* Since (2)  $\Rightarrow$  (3) and (3)  $\Rightarrow$  (1) are clear, we show (1)  $\Rightarrow$  (2). Let  $f$  be a monotone derivation. Since  $x \wedge y \leq x, y$ , we get  $f(x \wedge y) \leq f(x), f(y)$  and  $f(x \wedge y) \leq f(x) \wedge y, x \wedge f(y)$  by  $f(x \wedge y) \leq x \wedge y \leq x, y$ . On the other hand, since  $f$  is the derivation, we have  $f(x \wedge y) = (f(x) \wedge y) \vee (x \wedge f(y)) \geq f(x) \wedge y, x \wedge f(y)$ . This means that  $f(x \wedge y) = f(x) \wedge y = x \wedge f(y)$ . □

**Corollary 2.4.** *If  $L$  has a maximum element 1 and  $d$  is a derivation, then the following conditions are equivalent.*

- (1)  $d$  is monotone;
- (2)  $d(x) = x \wedge d(1)$  for all  $x \in L$ ;
- (3)  $d(x) \leq d(1)$  for all  $x \in L$ .

*Proof.* Since (1)  $\Rightarrow$  (3) and (2)  $\Rightarrow$  (1) are clear, we only show the case (3)  $\Rightarrow$  (2). Let  $d(x) \leq d(1)$ . Since  $d$  is the derivation, we have  $d(x) \leq x$  and thus  $d(x) \leq x \wedge d(1)$ . This implies  $d(x) = d(x \wedge 1) = (d(x) \wedge 1) \vee (x \wedge d(1)) = d(x) \vee (x \wedge d(1)) = x \wedge d(1)$ . □

**Corollary 2.5.** *If  $d$  is a monotone derivation of  $L$ , then  $d(d(x) \vee d(y)) = d(x) \vee d(y)$  for all  $x, y \in L$ .*

*Proof.* The proof follows from  $d(d(x) \vee d(y)) = d(d(x)) \vee d(d(y)) \leq d(d(x) \vee d(y)) \leq d(x) \vee d(y)$ .  $\square$

Unfortunately, the converse of the result above does not hold, namely,  $d$  may not be monotone even if  $d(d(x) \vee d(y)) = d(x) \vee d(y)$  holds. We have a counter example. Let  $L = \{0, a, b, 1\}$  with  $0 < a < b < 1$ . If we define  $d : L \rightarrow L$  by  $d(0) = d(1) = 0, d(a) = d(b) = b$ , then it is easy to show that  $d$  is a derivation and  $d(d(x) \vee d(y)) = d(x) \vee d(y)$ , but  $d$  is not monotone.

*Remark 2.6.* A map  $f : L \rightarrow L$  for a lattice  $L$  is called an *interior operator* if

- (io1)  $x \leq y \Rightarrow f(x) \leq f(y)$
- (io2)  $f(x) \leq x$
- (io3)  $f(f(x)) = f(x)$

It follows from our result above that a monotone derivation is an interior operator.

*Remark 2.7.* Similar results to our Theorem 2.2 are already proved in [7] as Theorem 3.19 and Theorem 3.21.

**Theorem 3.19.** Let  $L$  be a modular lattice and  $d$  be a derivation of  $L$ . Then the following conditions are equivalent:

- (1)  $d$  is monotone;
- (2)  $d(x \wedge y) = d(x) \wedge d(y)$ ;
- (3) If  $d(x) = x$ , then  $d(x \vee y) = d(x) \vee d(y)$ ,

where a lattice  $L$  is called *modular* if

$$x \leq z \Rightarrow x \vee (y \wedge z) = (x \vee y) \wedge z \quad (\text{for all } x, y, z \in L).$$

**Theorem 3.21.** Let  $L$  be a distributive lattice and  $d$  be a derivation of it. Then the following conditions are equivalent:

- (1)  $d$  is monotone.
- (2)  $d(x \wedge y) = d(x) \wedge d(y)$ .
- (3)  $d(x \vee y) = d(x) \vee d(y)$ .

Our results are stronger than those of above, because our results say that monotonicity is equivalent to the condition (2)  $d(x \wedge y) = d(x) \wedge d(y)$  for all lattices  $L$ , namely, we do not assume modularity nor distributivity to get such results.

Moreover, we obtain a following identity condition instead of (3) If  $d(x) = x$ , then  $d(x \vee y) = d(x) \vee d(y)$  in Theorem 3.19 in [7].

**Theorem 2.8.** *Let  $L$  be a modular lattice and  $d$  be a derivation. Then  $d$  is monotone  $\Leftrightarrow d(d(x) \vee y) = d(x) \vee d(y)$  ( $\forall x, y \in L$ )*

*Proof.* Suppose that  $d$  is monotone. Then we have

$$\begin{aligned} d(y) &= d(y \wedge (d(x) \vee y)) \\ &= \{d(y) \wedge (d(x) \vee y)\} \vee \{y \wedge d(d(x) \vee y)\} \\ &= d(y) \vee \{y \wedge d(d(x) \vee y)\} \quad (d(y) \leq y \leq d(x) \vee y) \\ &= (d(y) \vee d(d(x) \vee y)) \wedge y \quad (\text{modularity}) \\ &= y \wedge d(d(x) \vee y) \end{aligned}$$

and thus

$$\begin{aligned} d(x) \vee d(y) &= d(x) \vee \{y \wedge d(d(x) \vee y)\} \\ &= (d(x) \vee y) \wedge d(d(x) \vee y) \quad (\text{modularity}) \\ &= d(d(x) \vee y). \end{aligned}$$

Conversely, suppose  $d(d(x) \vee y) = d(x) \vee d(y)$ . If  $x \leq y$ , since  $d(x) \leq x \leq y$ , then we have  $d(x) \vee y = y$  and  $d(y) = d(d(x) \vee y) = d(x) \vee d(y)$ . Therefore  $d(x) \leq d(y)$  and  $d$  is monotone.  $\square$

Moreover we prove the converse.

**Theorem 2.9.** *A lattice  $L$  in which any derivation  $d$  satisfies the identity*

$$d(d(x) \vee y) = d(x) \vee d(y) \quad (\forall x, y \in L)$$

*is a modular lattice.*

*Proof.* For every  $z \in L$ , if we consider a map  $d_z(x) = x \wedge z$  then it is a monotone derivation. By assumption, the map  $d_z$  satisfies

$$d_z(d_z(x) \vee y) = d_z(x) \vee d_z(y) \quad (\forall x, y \in L)$$

and hence  $((x \wedge z) \vee y) \wedge z = (x \wedge z) \vee (y \wedge z)$ . This implies that if  $x \leq z$  then  $(x \vee y) \wedge z = x \vee (y \wedge z)$ . Therefore  $L$  is the modular lattice.  $\square$

We also have a similar result about distributive lattices.

**Theorem 2.10.** *Let  $L$  be a distributive lattice and  $d$  be a derivation. Then we have*

$$d \text{ is monotone} \Leftrightarrow d(x \vee y) = d(x) \vee d(y) \quad (\forall x, y \in L).$$

Conversely,

**Theorem 2.11.** *A lattice  $L$  in which any derivation  $d$  satisfies the identity*

$$d(x \vee y) = d(x) \vee d(y) \quad (\forall x, y \in L),$$

is a distributive lattice.

The above results provide characterization theorems of modular lattices and of distributive lattices in terms of derivations.

*Remark 2.12.* If  $d$  is a monotone derivation then a subset

$$\text{Fix}_d(L) = \{x \in L \mid d(x) = x\}$$

of  $L$  is an *ideal* of  $L$ , that is,  $\text{Fix}_d(L)$  satisfies the conditions

- (I1)  $0 \in \text{Fix}_d(L)$
- (I2)  $x \in \text{Fix}_d(L), y \leq x \Rightarrow y \in \text{Fix}_d(L)$
- (I3)  $x, y \in \text{Fix}_d(L) \Rightarrow x \vee y \in \text{Fix}_d(L)$ .

In the case of  $d$  being monotone, we have a following result.

**Theorem 2.13.** *If  $d$  is a monotone derivation of a lattice  $L$ , then  $\text{Fix}_d(L)$  is a lattice.*

*Proof.* For all  $x, y \in \text{Fix}_d(L)$ , since  $d$  is monotone, we have  $d(x \wedge y) = d(x) \wedge d(y) = x \wedge y$  and hence  $x \wedge y \in \text{Fix}_d(L)$ .  $\square$

*Remark 2.14.* We note that  $(\text{Fix}_d(L), \wedge, \vee)$  is a lattice for a monotone derivation  $d$ , but it is not always a sublattice of  $L$  if  $L$  has the maximum element 1. Because, if  $1 \in L$ , then  $(\text{Fix}_d(L), \wedge, \vee, 0, d(1))$  is also a lattice, however  $d(1) = 1$  does not hold in general.

**Corollary 2.15.** *If  $L$  is a bounded distributive lattice and  $d$  is a monotone derivation of  $L$ , then the quotient lattice  $L/\ker(d)$  is isomorphic to the lattice  $\text{Fix}_d(L)$ , that is,*

$$L/\ker(d) \cong \text{Fix}_d(L).$$

*Proof.* Let  $L$  be a bounded distributive lattice and  $d$  be a monotone derivation. Since  $d$  is monotone,  $d(z) = z \wedge d(1)$  for all  $z \in L$ . It follows that  $d(x \vee y) = (x \vee y) \wedge d(1) = (x \wedge d(1)) \vee (y \wedge d(1)) = d(x) \vee d(y)$ . This means that a map  $f : L \rightarrow \text{Fix}_d(L)$  defined by  $f(x) = d(x)$  for all  $x \in L$  is a surjective homomorphism. It follows from the homomorphism theorem of lattices that  $L/\ker(f) \cong \text{Fix}_d(L)$  and  $\ker(f) = \ker(d)$ , where  $x/\ker(d) = y/\ker(d)$  is defined by  $d(x) = d(y)$  for all  $x, y \in L$ . Therefore, we have  $L/\ker(d) \cong \text{Fix}_d(L)$ .  $\square$

### 3. OTHER DERIVATIONS

Some types of derivations, such as *generalized derivation*, *generalized  $(f, g)$ -derivation* and  *$f$ -derivation*, are defined and properties of them are considered in [1, 2, 3]. For instance, a map  $D : L \rightarrow L$  is called a

generalized derivation in [1] if it satisfies the condition: For a derivation  $d$ ,

$$D(x \wedge y) = (D(x) \wedge y) \vee (x \wedge d(y))$$

We get basic results about a generalized derivation  $D$  without difficulty.

**Proposition 3.1** (cf. Proposition 3.4, 3.9 [1]). *Let  $d$  be a derivation and  $D$  be a generalized derivation. Then we have*

- (1)  $d(x) \leq D(x) \leq x$ ;
- (1)  $D(D(x)) = D(x)$ ;
- (1)  $D(x) \wedge D(y) \leq D(x \wedge y)$ ;
- (1)  $D(x) \wedge D(y) = D(D(x) \wedge D(y))$ ;
- (1)  $D(x) = d(x) \vee (x \wedge D(1))$ .

We also have a new result about a generalized derivation  $D$ .

**Proposition 3.2.** *Let  $d$  be a derivation and  $D$  be a generalized derivation. Then we have  $D \circ d = d \leq d \circ D$*

*Proof.* Since

$$\begin{aligned} (D \circ d)(x) &= D(d(x)) \\ &= D(x \wedge d(x)) \\ &= (D(x) \wedge d(x)) \vee (x \wedge d(d(x))) \\ &= d(x) \vee (x \wedge d(x)) \\ &= d(x), \end{aligned}$$

we get  $D \circ d = d$ .

For  $d \circ D$ , we have  $d(D(x)) = d(x \wedge D(x)) = (d(x) \wedge D(x)) \vee (x \wedge d(D(x))) = d(x) \vee d(D(x)) \geq d(x)$  and hence  $d \leq d \circ D$ .  $\square$

It follows from our result that a characterization theorem about monotone generalized derivations can be proved similarly.

**Proposition 3.3** (Proposition 3.12 [1]). *For a generalized derivation  $D$ , the following conditions are equivalent to each other:*

- (1)  $D$  is monotone.
- (1)  $D(x \wedge y) = D(x) \wedge D(y)$ .
- (1)  $D(x) \vee D(y) \leq D(x \vee y)$ .
- (1)  $D(x) = x \wedge D(1)$  if  $L$  has a maximum element 1.

**Proposition 3.4.** *If  $L$  has a maximum element 1, then any generalized derivation  $D$  has a following form*

$$D(x) = (D(1) \wedge x) \vee d(x).$$

**Corollary 3.5.**  $D(1) = 1 \Leftrightarrow D = id_L$

**Lemma 3.6.** *If  $L$  has a maximum element  $1$  and  $d(x) \leq D(1)$  for all  $x \in L$ , then*

$$D(x) = x \wedge D(1).$$

In this case, the generalized derivation  $D$  is monotone. Conversely, if  $D$  is monotone then  $d(x) \leq D(1)$  for all  $x \in L$ . Therefore, we have another characterization of monotone generalized derivations.

**Theorem 3.7.** *For any generalized derivation  $D$ ,*

$$D \text{ is monotone} \Leftrightarrow d(x) \leq D(1) \quad (\forall x \in L).$$

**Corollary 3.8.** *If  $d$  is monotone, then so  $D$  is.*

*Proof.* We assume that  $d$  is monotone. Since  $d(x) = x \wedge d(1)$  ( $\forall x \in L$ ), we have  $d(x) = x \wedge d(1) \leq x \wedge D(1) \leq D(1)$  and thus  $D$  is monotone.  $\square$

We may ask whether the converse holds, that is, if a generalized derivation  $D$  is monotone then so  $d$  is ?

Unfortunately, this does not hold by the following example.

**Example 3** Let  $L = \{0, a, b, 1\}$ , ( $0 < a < b < 1$ ) and  $d, D : L \rightarrow L$  be maps defined by

$$d(x) = \begin{cases} 0 & (x = 0, 1) \\ a & (x = a, b) \end{cases}$$

$$D(x) = \begin{cases} x & (x = 0, a, b) \\ b & (x = 1) \end{cases}$$

It is easy to show that  $d$  is a derivation and  $D$  is a generalized derivation. Moreover  $D$  is monotone. However, it is obvious that  $d$  is not monotone.

In the previous section, we provide characterization theorems of modular lattices and of distributive lattices in terms of derivations. We also have similar results about generalized derivations.

**Theorem 3.9.** *Let  $L$  be a modular lattice and  $D$  be a generalized derivation. Then,  $D$  is monotone if and only if  $D(D(x) \vee y) = D(x) \vee D(y)$ .*

*Proof.* Suppose that  $D$  is monotone. Since

$$\begin{aligned} D(y) &= D((D(x) \vee y) \wedge y) \\ &= \{D(D(x) \vee y) \wedge y\} \vee \{(D(x) \vee y) \wedge d(y)\} \\ &= (D(D(x) \vee y) \wedge y) \vee d(y) \\ &= y \wedge D(D(x) \vee y), \end{aligned}$$

we have

$$\begin{aligned}
D(x) \vee D(y) &= D(x) \vee (y \wedge D(D(x) \vee y)) \\
&= D(D(x)) \vee (y \wedge D(D(x) \vee y)) \\
&= (D(x) \vee y) \wedge D(D(x) \vee y) \quad (\text{modularity}) \\
&= D(D(x) \vee y).
\end{aligned}$$

Conversely, suppose  $D(D(x) \vee y) = D(x) \vee D(y)$  for all  $x, y \in L$ . If  $x \leq y$ , since  $D(x) \leq x \leq y$ , then we have  $D(x) \vee y = y$  and  $D(y) = D(D(x) \vee y) = D(x) \vee D(y)$ . Therefore  $D(x) \leq D(y)$  and  $D$  is monotone.  $\square$

**Theorem 3.10.** *A lattice  $L$  in which any generalized derivation  $D$  satisfies the identity*

$$D(D(x) \vee y) = D(x) \vee D(y) \quad (\forall x, y \in L)$$

*is a modular lattice.*

*Proof.* For every  $z \in L$ , if we define maps  $d_z$  and  $D_z$  by  $d_z(x) = x \wedge z = D_z(x)$  for all  $x \in L$ . It is clear that  $d_z$  is a derivation and  $D_z$  is also a generalized derivation. Since  $D_z$  is monotone, it follows from assumption that  $D_z(D_z(x) \vee y) = D_z(x) \vee D_z(y)$  and thus  $((x \wedge z) \vee y) \wedge z = (x \wedge z) \vee (y \wedge z)$ . This implies that if  $x \leq z$  then  $(x \vee y) \wedge z = x \vee (y \wedge z)$ . Therefore  $L$  is the modular lattice.  $\square$

We also have a similar result about distributive lattices.

**Theorem 3.11** (Theorem 3.14 [1]). *Let  $L$  be a distributive lattice and  $D$  be a generalized derivation. Then we have*

$$D \text{ is monotone} \Leftrightarrow D(x \vee y) = D(x) \vee D(y) \quad (\forall x, y \in L).$$

Conversely,

**Theorem 3.12.** *A lattice  $L$  in which any generalized derivation  $D$  satisfies the identity*

$$D(x \vee y) = D(x) \vee D(y) \quad (\forall x, y \in L)$$

*is a distributive lattice.*

The above results provide characterization theorems of modular lattices and of distributive lattices in terms of generalized derivations.

We also consider another type of derivation, *f-derivation*, according to [3]. A map  $d : L \rightarrow L$  is called an *f-derivation* if there exists a map  $f : L \rightarrow L$  such that

$$d(x \wedge y) = (d(x) \wedge f(y)) \vee (f(x) \wedge d(y)) \quad (\forall x, y \in L).$$

It is clear that if  $f = id_L$  then an  $f$ -derivation is the same as the derivation defined in the previous section.

As basic results about  $f$ -derivations, we have

**Proposition 3.13** (Proposition 1,2 [3]). *Let  $d : L \rightarrow L$  be an  $f$ -derivation. Then, for all  $x, y \in L$ ,*

- (1)  $d(x) \leq f(x)$ ;
- (2)  $d(x) \wedge d(y) \leq d(x \wedge y) \leq d(x) \vee d(y)$ ;
- (3)  $f(x) \leq d(1), f(1) = 1 \Rightarrow d(x) = f(x)$ ;
- (4)  $d(1) = 1 \Rightarrow d = f$ , hence  $d$  is monotone.

We also have similar results about monotone  $f$ -derivations.

**Proposition 3.14** (cf. Theorem 1 [3]). *For an  $f$ -derivation  $d$ , the following conditions are equivalent:*

- (1)  $d$  is monotone;
- (2)  $d(x) \vee d(y) \leq d(x \vee y) \quad (\forall x, y \in L)$ ;
- (3)  $d(x \wedge y) = d(x) \wedge d(y) \quad (\forall x, y \in L)$ ;
- (4)  $d(x) = f(x) \wedge d(x \vee y) \quad (\forall x, y \in L)$ .

*Proof.* We only show that (1) is equivalent to (4). Suppose that  $d$  is monotone. Since  $d(x) \leq f(x)$  and  $d(x) \leq d(x \vee y)$ , we get  $d(x) \leq f(x) \wedge d(x \vee y)$ . On the other hand, since  $d(x) \leq d(x \vee y) \leq f(x \vee y)$ , we have

$$\begin{aligned} d(x) &= d(x \wedge (x \vee y)) \\ &= (d(x) \wedge f(x \vee y)) \vee (f(x) \wedge d(x \vee y)) \quad (\because d \text{ is an } f\text{-derivation}) \\ &= d(x) \vee (f(x) \wedge d(x \vee y)) \quad (\because d(x) \leq f(x \vee y)). \end{aligned}$$

This means that  $f(x) \wedge d(x \vee y) \leq d(x)$ . Therefore, we get  $d(x) = f(x) \wedge d(x \vee y)$ .

Conversely, we assume that  $d(x) = f(x) \wedge d(x \vee y) \quad (\forall x, y \in L)$ . If  $x \leq y$ , then  $d(x) = f(x) \wedge d(x \vee y) = f(x) \wedge d(y) \leq d(y)$ . Hence  $d$  is monotone. □

We note that the result above was already proved in [3] as Theorem 1 under the conditions  $f(1) = 1$  and  $f(x \wedge y) = f(x) \wedge f(y)$  for all  $x, y \in L$ . Our result shows that the conditions are redundant. Moreover, our result implies that the modularity condition is also redundant in the Theorem 2 (a) in [3], where it said that

Theorem 2. Let  $L$  be a modular lattice and  $d$  be an  $f$ -derivation on  $L$ .

- (a)  $d$  is a monotone  $f$ -derivation if and only if  $d(x \wedge y) = dx \wedge dy$ .

In cases of modular lattices and of distributive lattices, we have following results.

**Theorem 3.15.** *Let  $L$  be a modular lattice and  $d : L \rightarrow L$  be an  $f$ -derivation. Then,  $d$  is monotone if and only if  $d(x) \vee d(y) = (d(x) \vee f(y)) \wedge d(x \vee y)$  for all  $x, y \in L$ .*

*Proof.* Suppose that  $d$  is monotone. Since  $d(y) = f(y) \wedge d(x \vee y)$ , we have  $d(x) \vee d(y) = d(x) \vee (f(y) \wedge d(x \vee y)) = (d(x) \vee f(y)) \wedge d(x \vee y)$  by modularity.

Conversely, assume that  $d(x) \vee d(y) = (d(x) \vee f(y)) \wedge d(x \vee y)$  for all  $x, y \in L$ . If  $x \leq y$  then  $d(x) \leq d(x) \vee d(y) = (d(x) \vee f(y)) \wedge d(x \vee y) \leq d(y)$  and thus  $d$  is monotone.  $\square$

**Corollary 3.16.** *Let  $L$  be a modular lattice and  $d : L \rightarrow L$  be a derivation. Then,  $d$  is monotone if and only if  $d(d(x) \vee y) = d(x) \vee d(y)$ .*

*Proof.* For an  $f$ -derivation  $d$  of a modular lattice  $L$ , if we take  $f = id_L$ , then  $d$  is a derivation of  $L$  and thus  $d \circ d = d$  and  $d(x) \leq x$  for all  $x \in L$ . It follows from the above that  $d(x) \vee d(y) = (d(x) \vee y) \wedge d(x \vee y)$  for all  $x, y \in L$ . By use of these facts, if  $d$  is monotone, then we have  $d(x) \vee d(y) = d(d(x)) \vee d(y) = (d(d(x)) \vee y) \wedge d(d(x) \vee y) = (d(x) \vee y) \wedge d(d(x) \vee y) = d(d(x) \vee y)$ . The converse is obvious.  $\square$

For the case of distributive lattices, we also have a following result.

**Theorem 3.17.** *Let  $L$  be a distributive lattice and  $d$  be an  $f$ -derivation. Then,  $d$  is monotone if and only if  $d(x) \vee d(y) = (f(x) \vee f(y)) \wedge d(x \vee y)$  for all  $x, y \in L$ .*

*Proof.* Let  $d$  be a monotone  $f$ -derivation. Since  $d(x) = f(x) \wedge d(x \vee y)$  and  $L$  is the distributive lattice, we have  $d(x) \vee d(y) = (f(x) \wedge d(x \vee y)) \vee (f(y) \wedge d(x \vee y)) = (f(x) \vee f(y)) \wedge d(x \vee y)$ .

Conversely, suppose that  $d(x) \vee d(y) = (f(x) \vee f(y)) \wedge d(x \vee y)$  for all  $x, y \in L$ . If  $x \leq y$ , then  $d(x) \leq d(x) \vee d(y) = (f(x) \vee f(y)) \wedge d(x \vee y) \leq d(y)$ . Therefore  $d$  is monotone.  $\square$

**Corollary 3.18.** *Let  $L$  be a distributive lattice and  $d$  be a derivation. Then,  $d$  is monotone if and only if  $d(x) \vee d(y) = d(x \vee y)$  for all  $x, y \in L$ .*

*Remark 3.19.* The following result was proved as theorem 4 which was one of the main results of [3].

Theorem 4. Let  $L$  be a lattice. If there exists an  $f$ -derivation  $d$  on  $L$  such that  $d(x \vee y) = d(x) \vee d(y)$  for all  $x, y \in L$  and  $f$  is an epimorphism, then  $L$  is a distributive lattice.

Unfortunately, this is not true, because we have a following counter example. Let  $L = N_5 = \{0, a, b, c, 1\}$ , ( $0 < a < 1, 0 < b < c < 1$ ) and  $f = d = id_L$ . Then it is trivial that  $d$  and  $f$  satisfy the assumption of the theorem, but the lattice  $L$  is neither distributive nor modular.

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#### Mayuka F. Kawaguchi

Graduate School of Information Science and Technology, Hokkaido University,  
P.O.Box 060-0814 Sapporo, Japan  
Email: mayuka@ist.hokudai.ac.jp

#### Michiro Kondo

Department of Mathematics, School of System Design and Technology, Tokyo Denki  
University, P.O.Box 120-8551 Tokyo, Japan  
Email: mkondo@mail.dendai.ac.jp

SOME PROPERTIES ON DERIVATIONS OF LATTICES

M. F. KAWAGUCHI AND M. KONDO

برخی خواص مشتقات شبکه‌ها

<sup>۱</sup>مایوکا کاواگوچی و <sup>۲</sup>میشیرو کوندو

<sup>۱</sup> دانشکده تحصیلات تکمیلی علوم و فناوری اطلاعات، دانشگاه هوکایدو، ساپورو، ژاپن

<sup>۲</sup> گروه ریاضی، دانشکده طراحی سیستم و فناوری، دانشگاه توکیو دنکی، توکیو، ژاپن

در این مقاله برخی خواص مشتقات شبکه‌ها را بررسی کرده و نشان می‌دهیم (i) مشتق  $d$  از شبکه  $L$  با عنصر ماکسیم  $۱$ ، یکنوا است اگر و تنها اگر  $d(x) \leq d(1)$  برای هر  $x \in L$ ، (ii) مشتق یکنوای  $d$  توسط  $d(x) = x \wedge d(1)$  مشخص (توصیف) می‌شود و (iii) برخی قضایای رده بندی ساده برای شبکه‌های پیمانهای و شبکه‌های توزیع‌پذیر (پخشی)، با استفاده از مشتقات شبکه‌ها بیان خواهد شد. همچنین نشان خواهیم داد که برای شبکه توزیع‌پذیر  $L$  و مشتق یکنوای  $d$  از این شبکه، مجموعه  $\text{Fix}_d(L)$  متشکل از تمام نقاط ثابت  $d$  با شبکه  $L/\ker(d)$  یکرخت می‌باشد. به‌علاوه، مثال نقضی برای قضیه ۴ که از مرجع [۳]، ارائه خواهیم کرد.

کلمات کلیدی: مشتق، حافظ ترتیب، شبکه پیمانهای، شبکه توزیع‌پذیر.